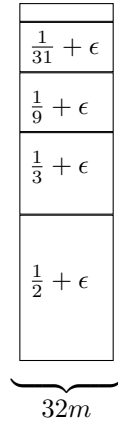


## Algorithms Theory, Solution for Assignment 7

[http://lak.informatik.uni-freiburg.de/lak\\_teaching/ws09\\_10/algo0910.php](http://lak.informatik.uni-freiburg.de/lak_teaching/ws09_10/algo0910.php)

### Exercise 7.1 - Bin packing

- An optimal packing yields  $\text{OPT}(I) = 32m$  bins.



- The strategy First Fit Decreasing first sorts the sequence  $I$ :

$$\underbrace{\frac{1}{2} + \epsilon, \dots, \frac{1}{2} + \epsilon}_{32m} \quad \underbrace{\frac{1}{3} + \epsilon, \dots, \frac{1}{3} + \epsilon}_{32m} \quad \underbrace{\frac{1}{9} + \epsilon, \dots, \frac{1}{9} + \epsilon}_{32m} \quad \underbrace{\frac{1}{31} + \epsilon, \dots, \frac{1}{31} + \epsilon}_{32m}$$

Then, it applies First Fit which results in the optimal packing showed above using  $32m$  bins.

### Exercise 7.2 - Matrix-chain multiplication

Find the optimal parenthesization of a matrix-chain product whose sequence  $P$  of dimensions is  $\langle 44, 43, 3, 29, 35, 19 \rangle$ . Specify all values  $m[i, j]$  and  $s[i, j]$ . Finally, provide the optimal parenthesization for  $A_1 \cdot A_2 \cdot A_3 \cdot A_4 \cdot A_5$ .

$$P = \langle p_0, p_1, p_2, p_3, p_4, p_5 \rangle$$

- l=1:

$$- m[i, i] = 0 \quad i \in \{1, 2, 3, 4, 5\}$$

- l=2:

$$- m[1, 2] = \min_{k=1} (m[1, k] + m[k + 1, 2] + p_0 p_1 p_2) = 5676$$

$$- m[2, 3] = \min_{k=2} (m[2, k] + m[k + 1, 3] + p_1 p_2 p_3) = 3741$$

$$- m[3, 4] = \min_{k=3} (m[3, k] + m[k + 1, 4] + p_2 p_3 p_4) = 3045$$

$$- m[4, 5] = \min_{k=4} (m[4, k] + m[k + 1, 5] + p_3 p_4 p_5) = 19285$$

- l=3:

$$- m[1, 3] = \min_{k=1,2} (m[1, k] + m[k + 1, 3] + p_0 p_k p_3)$$

$$m[1, 3] = \min \begin{cases} m[1, 1] + m[2, 3] + p_0 p_1 p_3 = 0 + 3741 + 54868 = 58609 \\ m[1, 2] + m[3, 3] + p_0 p_2 p_3 = 5676 + 0 + 3828 = 9504 \end{cases} = 9504$$

$$s[1, 3] = 2$$

$$\begin{aligned}
& - m[2, 4] = \min_{k=2,3}(m[2, k] + m[k + 1, 4] + p_1 p_k p_4) \\
& m[2, 4] = \min \begin{cases} m[2, 2] + m[3, 4] + p_1 p_2 p_4 = 0 + 3045 + 4515 = 7560 \\ m[2, 3] + m[4, 4] + p_1 p_3 p_4 = 3741 + 0 + 43645 = 47386 \end{cases} = 7560 \\
& s[2, 4] = 2 \\
& - m[3, 5] = \min_{k=3,4}(m[3, k] + m[k + 1, 5] + p_2 p_k p_5) \\
& m[3, 5] = \min \begin{cases} m[3, 3] + m[4, 5] + p_2 p_3 p_5 = 0 + 19285 + 1653 = 20938 \\ m[3, 4] + m[5, 5] + p_2 p_4 p_5 = 3045 + 0 + 1995 = 5040 \end{cases} = 5040 \\
& s[3, 5] = 4
\end{aligned}$$

• l=4:

$$\begin{aligned}
& - m[1, 4] = \min_{k=1,2,3}(m[1, k] + m[k + 1, 4] + p_0 p_k p_4) \\
& m[1, 4] = \min \begin{cases} m[1, 1] + m[2, 4] + p_0 p_1 p_4 = 0 + 7560 + 66220 = 73780 \\ m[1, 2] + m[3, 4] + p_0 p_2 p_4 = 5676 + 3045 + 4620 = 13341 \\ m[1, 3] + m[4, 4] + p_0 p_3 p_4 = 9504 + 0 + 44660 = 54164 \end{cases} = 13341 \\
& s[1, 4] = 2 \\
& - m[2, 5] = \min_{k=2,3,4}(m[2, k] + m[k + 1, 5] + p_1 p_k p_5) \\
& m[2, 5] = \min \begin{cases} m[2, 2] + m[3, 5] + p_1 p_2 p_5 = 0 + 5040 + 2451 = 7491 \\ m[2, 3] + m[4, 5] + p_1 p_3 p_5 = 3741 + 19285 + 23693 = 46719 \\ m[2, 4] + m[5, 5] + p_1 p_4 p_5 = 7560 + 0 + 28595 = 36155 \end{cases} = 7491 \\
& s[2, 5] = 2
\end{aligned}$$

• l=4:

$$\begin{aligned}
& - m[1, 5] = \min_{k=1,2,3,4}(m[1, k] + m[k + 1, 5] + p_0 p_k p_5) \\
& m[1, 5] = \min \begin{cases} m[1, 1] + m[2, 5] + p_0 p_1 p_5 = 0 + 7491 + 35948 = 43439 \\ m[1, 2] + m[3, 5] + p_0 p_2 p_5 = 5676 + 5040 + 2508 = 13224 \\ m[1, 3] + m[4, 5] + p_0 p_3 p_5 = 9504 + 19285 + 24244 = 53033 \\ m[1, 4] + m[5, 5] + p_0 p_4 p_5 = 13341 + 0 + 29260 = 42601 \end{cases} = 13224 \\
& s[1, 5] = 2
\end{aligned}$$

Resulting tables:

		$j$				
		1	2	3	4	5
$i$	1	0	5676	9504	13341	13224
	2		0	3741	7560	7491
	3			0	3045	5040
	4				0	19285
	5					0

		$j$				
		$s$	2	3	4	5
$i$	1	1	2	2	2	2
	2			2	2	2
	3				3	4
	4					4

Final parenthesization:  $(A_1 A_2)((A_3 A_4) A_5)$

### Exercise 7.3 - Optimal search trees

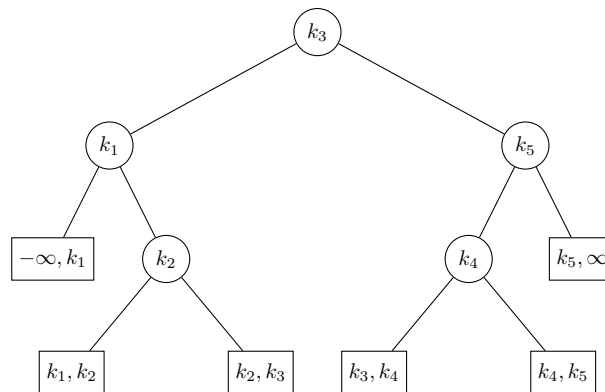
Applying the algorithm we obtain the following tables:

		$j$					
$W$		0	1	2	3	4	5
$i$	0	4	12	17	25	35	43
	1		3	8	16	26	34
	2			2	10	20	28
	3				1	11	19
	4					4	12
	5						3

		$j$					
$P$		0	1	2	3	4	5
$i$	0	0	12	25	47	71	98
	1		0	8	24	45	70
	2			0	10	30	50
	3				0	11	30
	4					0	12
	5						0

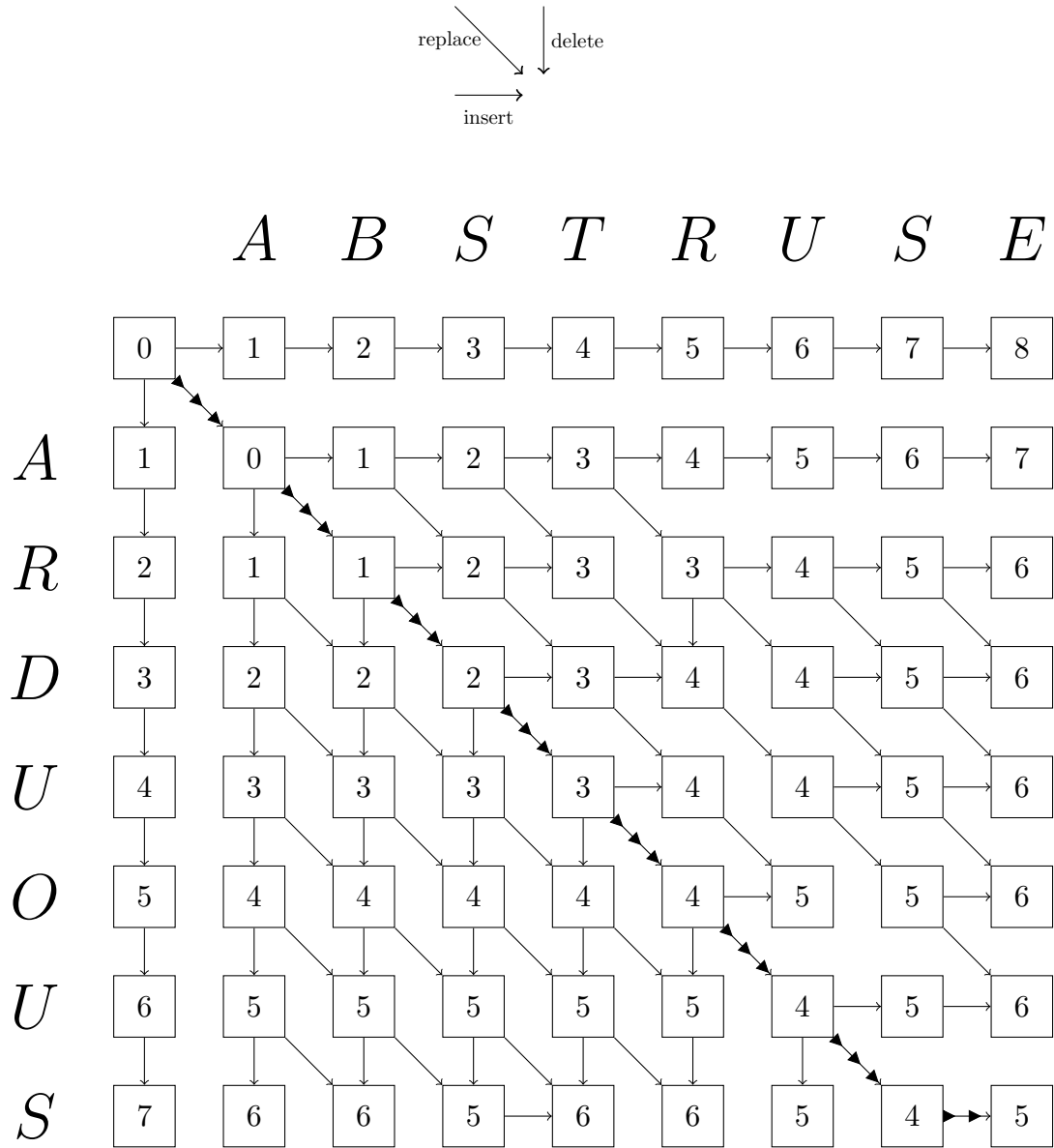
		$j$				
$r$		1	2	3	4	5
$i$	0	1	1	2	3	3
	1		2	3	3	4
	2			3	4	4
	3				4	5
	4					5

Optimal search tree:



### Exercise 7.4 - Edit Distance

- Trace graph:



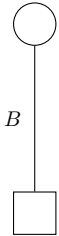
- Sequence:  $D(A, B) = 5$

Operation	Cost
replace A by A	0
replace R by B	1
replace D by S	1
replace U by T	1
replace O by R	1
replace U by U	0
replace S by S	0
insert E	1

### Exercise 7.5 - Text search (Ukkonen algorithm)

- $T_1$  :

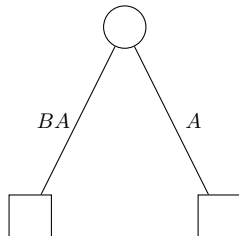
- rule 2



- $T_2$  :

- apply implicit extensions

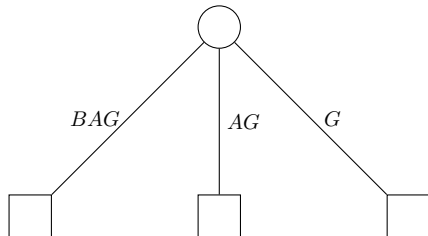
- ext 2, rule 2



- $T_3$  :

- apply implicit extensions

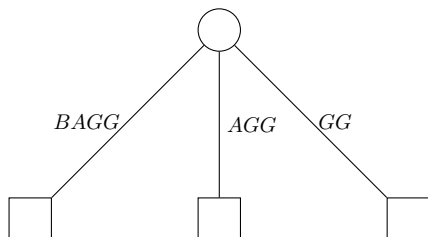
- ext 3, rule 2



- $T_4$  :

- apply implicit extensions

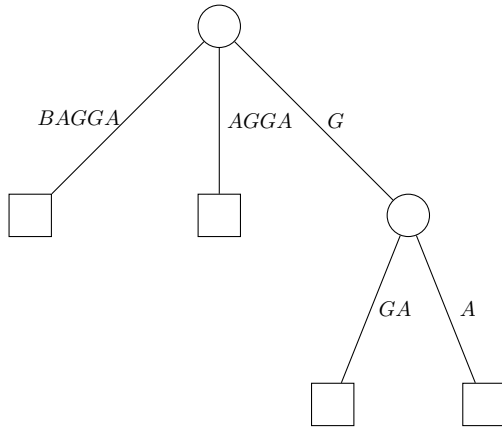
- ext 4, rule 3



- $T_5$  :

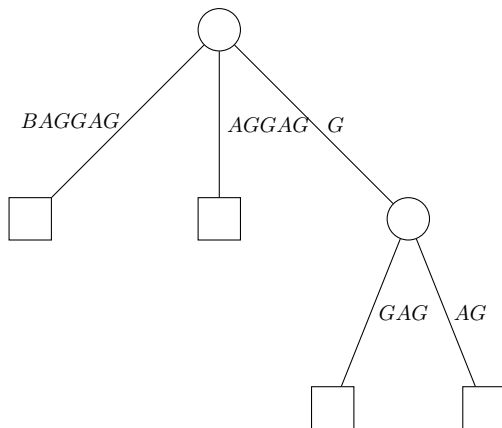
- apply implicit extensions

- ext 4, rule 2
- ext 5, rule 3



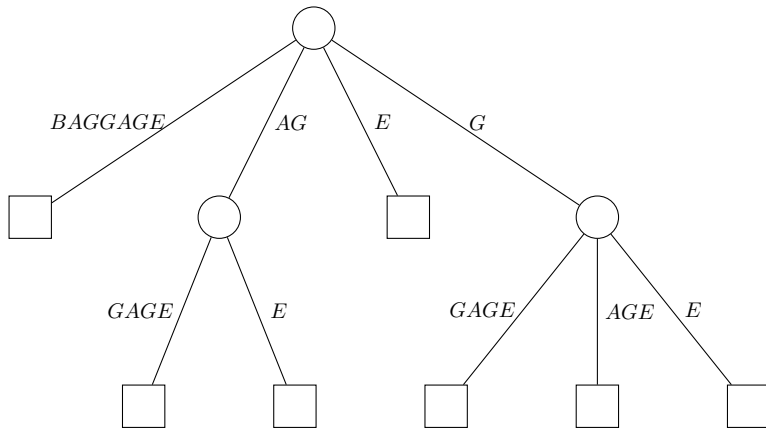
•  $T_6$  :

- apply implicit extensions
- ext 5, rule 3, phase ended



•  $T_7$  :

- apply implicit extensions
- ext 5, rule 2
- ext 6, rule 2
- ext 7, rule 2



•  $\underline{T}$ : BAGGAGE

